

A First Look At Prolog

Outline

- Terms
- Using a Prolog language system
- Rules
- The two faces of Prolog
- Operators
- Lists
- Negation and failure
- What Prolog is good for

Terms

- Everything in Prolog is built from *terms*:
 - Prolog programs
 - The data manipulated by Prolog programs
- Three kinds of terms:
 - Constants: integers, real numbers, atoms
 - Variables
 - Compound terms

Constants

- Integer constants: **123**
- Real constants: **1.23**
- Atoms:
 - A lowercase letter followed by any number of additional letters, digits or underscores: **fred**
 - A sequence of non-alphanumeric characters: ***, ., =, @#\$**
 - Plus a few special atoms: **[]**

Atoms Are Not Variables

- An atom can look like an ML or Java variable:
 - **i**, **size**, **length**
- But an atom is not a variable; it is not bound to anything, never equal to anything else
- Think of atoms as being more like string constants: "**i**", "**size**", "**length**"

Variables

- Any name beginning with an uppercase letter or an underscore, followed by any number of additional letters, digits or underscores: **X, Child, Fred, _, _123**
- Most of the variables you write will start with an uppercase letter
- Those starting with an underscore, including **_**, get special treatment

Compound Terms

- An atom followed by a parenthesized, comma-separated list of one or more terms:
 **$x(y, z)$, $+ (1, 2)$, $. (1, [])$,
parent(adam, seth), **$x(Y, x(Y, Z))$****
- A compound term can look like an ML function call: **$f(x, y)$**
- Again, this is misleading
- Think of them as structured data

Terms

$\langle term \rangle ::= \langle constant \rangle \mid \langle variable \rangle \mid \langle compound-term \rangle$
 $\langle constant \rangle ::= \langle integer \rangle \mid \langle real\ number \rangle \mid \langle atom \rangle$
 $\langle compound-term \rangle ::= \langle atom \rangle (\langle termlist \rangle)$
 $\langle termlist \rangle ::= \langle term \rangle \mid \langle term \rangle , \langle termlist \rangle$

- All Prolog programs and data are built from such terms
- Later, we will see that, for instance, **+ (1, 2)** is usually written as **1+2**
- But these are not new kinds of terms, just abbreviations

Unification

- Pattern-matching using Prolog terms
- Two terms *unify* if there is some way of binding their variables that makes them identical
- For instance, **parent (adam, Child)** and **parent (adam, seth)** unify by binding the variable **Child** to the atom **seth**
- More details later: Chapter 20

The Prolog Database

- A Prolog language system maintains a collection of facts and rules of inference
- It is like an internal database that changes as the Prolog language system runs
- A Prolog program is just a set of data for this database
- The simplest kind of thing in the database is a *fact*: a term followed by a period

Example

```
parent (kim, holly) .  
parent (margaret, kim) .  
parent (margaret, kent) .  
parent (esther, margaret) .  
parent (herbert, margaret) .  
parent (herbert, jean) .
```

- A Prolog program of six facts
- Defining a *predicate* **parent** of *arity* 2
- We would naturally interpret these as facts about families: Kim is the parent of Holly and so on

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SWI-Prolog

```
Welcome to SWI-Prolog ...
```

```
For help, use ?- help(Topic) . or ?- apropos(Word) .
```

```
?-
```

- Prompting for a query with `?-`
- Normally interactive: get query, print result, repeat

The **consult** Predicate

```
?- consult(relations).  
% relations compiled 0.00 sec, 852 bytes  
true.  
  
?-
```

- Predefined predicate to read a program from a file into the database
- File **relations** (or **relations.pl**) contains our **parent** facts

Simple Queries

```
?- parent (margaret, kent) .  
true.  
  
?- parent (fred, pebbles) .  
false.  
  
?-
```

- A query asks the language system to prove something
- Some turn out to be **true**, some **false**
- (Some queries, like **consult**, are executed only for their side-effects)

Final Period

```
?- parent (margaret, kent)
|      .
true.

?-
```

- Queries can take multiple lines
- If you forget the final period, Prolog prompts for more input with |

Queries With Variables

```
?- parent (P, jean) .  
P = herbert.  
  
?- parent (P, esther) .  
false.
```

- Any term can appear as a query, including a term with variables
- The Prolog system shows the bindings necessary to prove the query

Flexibility

- Normally, variables can appear in any or all positions in a query:
 - **parent (Parent , jean)**
 - **parent (esther , Child)**
 - **parent (Parent , Child)**
 - **parent (Person , Person)**

Multiple Solutions

```
?- parent (Parent, Child) .  
Parent = kim,  
Child = holly .
```

- When the system finds a solution, it prints the binding it found
- If it could continue to search for additional solutions, it then prompts for input
- Hitting Enter makes it stop searching and print the final period...

Multiple Solutions

- ... entering a semicolon makes it continue the search
- As often as you do this, it will try to find another solution
- In this case, there is one for every fact in the database

```
?- parent (Parent, Child) .  
Parent = kim,  
Child = holly ;  
Parent = margaret,  
Child = kim ;  
Parent = margaret,  
Child = kent ;  
Parent = esther,  
Child = margaret ;  
Parent = herbert,  
Child = margaret ;  
Parent = herbert,  
Child = jean .
```

Conjunctions

```
?- parent(margaret, X), parent(X, holly).  
X = kim .
```

- A conjunctive query has a list of query terms separated by commas
- The Prolog system tries prove them all (using a single set of bindings)

```
?- parent (Parent, kim), parent (Grandparent, Parent) .  
Parent = margaret,  
Grandparent = esther ;  
Parent = margaret,  
Grandparent = herbert ;  
false.
```

```
?- parent (esther, Child),  
|   parent (Child, Grandchild),  
|   parent (Grandchild, GreatGrandchild) .  
Child = margaret,  
Grandchild = kim,  
GreatGrandchild = holly .
```

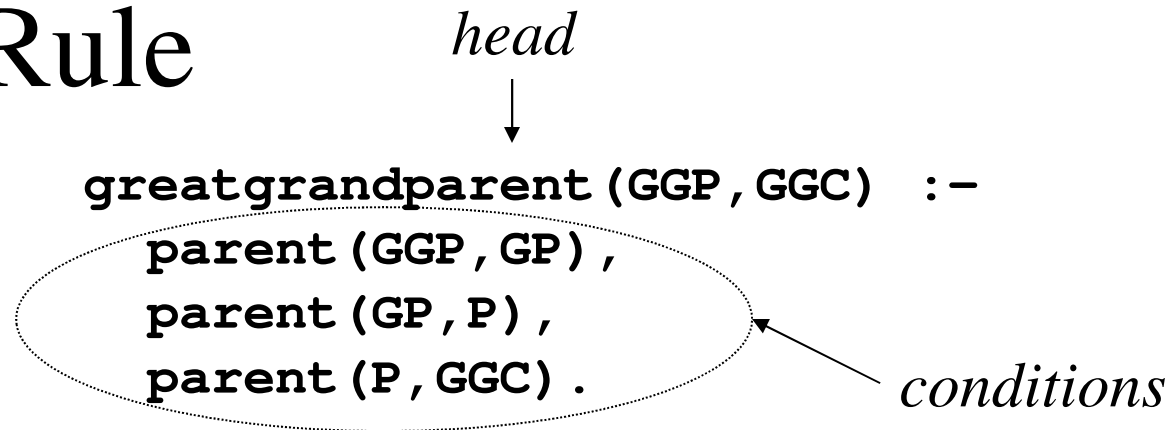
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The Need For Rules

- Previous example had a lengthy query for great-grandchildren of Esther
- It would be nicer to query directly:
greatgrandparent (esther, GGC)
- But we do not want to add separate facts of that form to the database
- The relation should follow from the **parent** relation already defined

A Rule



- A rule says how to prove something: to prove the head, prove the conditions
- To prove **greatgrandparent (GGP, GGC)**, find some **GP** and **P** for which you can prove **parent (GGP, GP)**, then **parent (GP, P)** and then finally **parent (P, GGC)**

A Program With The Rule

```
parent (kim, holly) .  
parent (margaret, kim) .  
parent (margaret, kent) .  
parent (esther, margaret) .  
parent (herbert, margaret) .  
parent (herbert, jean) .  
greatgrandparent (GGP, GGC) :-  
    parent (GGP, GP) , parent (GP, P) , parent (P, GGC) .
```

- A program consists of a list of *clauses*
- A clause is either a fact or a rule, and ends with a period

Example

```
?- greatgrandparent (esther, GreatGrandchild) .  
GreatGrandchild = holly .
```

- This shows the initial query and final result
- Internally, there are intermediate *goals*:
 - The first goal is the initial query
 - The next is what remains to be proved after transforming the first goal using one of the clauses (in this case, the greatgrandparent rule)
 - And so on, until nothing remains to be proved

1. **parent (kim,holly) .**
2. **parent (margaret,kim) .**
3. **parent (margaret,kent) .**
4. **parent (esther,margaret) .**
5. **parent (herbert,margaret) .**
6. **parent (herbert,jean) .**
7. **greatgrandparent (GGP,GGC) :-**
 parent (GGP,GP) , parent (GP,P) , parent (P,GGC) .

*We will see more
about Prolog's model
of execution in
Chapter 20*

greatgrandparent (esther,GreatGrandchild)

↓ Clause 7, binding **GGP** to **esther** and **GGC** to **GreatGrandChild**

parent (esther,GP) , parent (GP,P) , parent (P,GreatGrandchild)

↓ Clause 4, binding **GP** to **margaret**

parent (margaret,P) , parent (P,GreatGrandchild)

↓ Clause 2, binding **P** to **kim**

parent (kim,GreatGrandchild)

↓ Clause 1, binding **GreatGrandchild** to **holly**

Rules Using Other Rules

```
grandparent (GP, GC) :-  
    parent (GP, P), parent (P, GC) .
```

```
greatgrandparent (GGP, GGC) :-  
    grandparent (GGP, P), parent (P, GGC) .
```

- Same relation, defined indirectly
- Note that both clauses use a variable **P**
- The scope of the definition of a variable is the clause that contains it

Recursive Rules

```
ancestor(X, Y) :- parent(X, Y) .  
ancestor(X, Y) :-  
    parent(Z, Y) ,  
    ancestor(X, Z) .
```

- **X** is an ancestor of **Y** if:
 - Base case: **X** is a parent of **Y**
 - Recursive case: there is some **Z** such that **Z** is a parent of **Y**, and **X** is an ancestor of **Z**
- Prolog tries rules in the order you give them, so put base-case rules and facts first

```
?- ancestor(jean, jean) .  
false.
```

```
?- ancestor(kim, holly) .  
true .
```

```
?- ancestor(A, holly) .  
A = kim ;  
A = margaret ;  
A = esther ;  
A = herbert ;  
false.
```

Core Syntax Of Prolog

- You have seen the complete core syntax:

$$\begin{aligned} \langle clause \rangle & ::= \langle fact \rangle \mid \langle rule \rangle \\ \langle fact \rangle & ::= \langle term \rangle . \\ \langle rule \rangle & ::= \langle term \rangle :- \langle termlist \rangle . \\ \langle termlist \rangle & ::= \langle term \rangle \mid \langle term \rangle , \langle termlist \rangle \end{aligned}$$

- There is not much more syntax for Prolog than this: it is a very simple language
- Syntactically, that is!

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The Procedural Side

```
greatgrandparent (GGP, GGC) :-  
    parent (GGP, GP), parent (GP, P), parent (P, GGC) .
```

- A rule says how to prove something:
 - To prove `greatgrandparent (GGP, GGC)`, find some `GP` and `P` for which you can prove `parent (GGP, GP)`, then `parent (GP, P)` and then finally `parent (P, GGC)`
- A Prolog program specifies proof procedures for queries

The Declarative Side

- A rule is a logical assertion:
 - For all bindings of **GGP**, **GP**, **P**, and **GGC**, if **parent (GGP, GP)** and **parent (GP, P)** and **parent (P, GGC)**, then **greatgrandparent (GGP, GGC)**
- Just a formula – it doesn't say how to *do* anything – it just makes an assertion:

$$\forall GGP, GP, P, GGC . \text{parent}(GGP, GP) \wedge \text{parent}(GP, P) \wedge \text{parent}(P, GGC) \\ \Rightarrow \text{greatgrandparent}(GGP, GGC)$$

Declarative Languages

- Each piece of the program corresponds to a simple mathematical abstraction
 - Prolog clauses – formulas in first-order logic
 - ML fun definitions – functions
- Many people use *declarative* as the opposite of *imperative*, including both logic languages and functional languages

Declarative Advantages

- Imperative languages are doomed to subtle side-effects and interdependencies
- Simpler declarative semantics makes it easier to develop and maintain correct programs
- Higher-level, more like *automatic programming*: describe the problem and have the computer write the program

Prolog Has Both Aspects

- Partly declarative
 - A Prolog program has logical content
- Partly procedural
 - A Prolog program has procedural concerns: clause ordering, condition ordering, side-effecting predicates, etc.
- It is important to be aware of both

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Operators

- Prolog has some predefined operators (and the ability to define new ones)
- An operator is just a predicate for which a special abbreviated syntax is supported

The = Predicate

- The goal $= (\mathbf{X}, \mathbf{Y})$ succeeds if and only if \mathbf{X} and \mathbf{Y} can be unified:

```
?- = (parent (adam, seth) , parent (adam, X) ) .  
X = seth.
```

- Since = is an operator, it can be and usually is written like this:

```
?- parent (adam, seth) =parent (adam, X) .  
X = seth.
```

Arithmetic Operators

- Predicates $+$, $-$, $*$ and $/$ are operators too, with the usual precedence and associativity

```
?- X = +(1, *(2, 3)).  
X = 1+2*3.
```

```
?- X = 1+2*3.  
X = 1+2*3.
```

Prolog lets you use operator notation, and prints it out that way, but the underlying term is still $+(1, *(2, 3))$

Not Evaluated

```
?- +(X, Y) = 1+2*3.  
X = 1,  
Y = 2*3.  
  
?- 7 = 1+2*3.  
false.
```

- The term is still **+ (1, * (2, 3))**
- It is not evaluated
- There is a way to make Prolog evaluate such terms, but we won't need it yet

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Lists in Prolog

- A bit like ML lists
- The atom `[]` represents the empty list
- A predicate `.` corresponds to ML's `::` operator

ML expression	Prolog term
<code>[]</code>	<code>[]</code>
<code>1 :: []</code>	<code>. (1, [])</code>
<code>1 :: 2 :: 3 :: []</code>	<code>. (1, . (2, . (3, [])))</code>
No equivalent.	<code>. (1, . (parent (X, Y), []))</code>

List Notation

List notation	Term denoted
<code>[]</code>	<code>[]</code>
<code>[1]</code>	<code>. (1, [])</code>
<code>[1, 2, 3]</code>	<code>. (1, . (2, . (3, [])))</code>
<code>[1, parent (X, Y)]</code>	<code>. (1, . (parent (X, Y) , []))</code>

- ML-style notation for lists
- These are just abbreviations for the underlying term using the `.` Predicate
- Prolog usually displays lists in this notation

Example

```
?- X = .(1, .(2, .(3, []))) .  
X = [1, 2, 3] .
```

```
?- .(X, Y) = [1, 2, 3] .  
X = 1,  
Y = [2, 3] .
```

List Notation With Tail

List notation	Term denoted
$[1 \mid \mathbf{x}]$	$. (1, \mathbf{x})$
$[1, 2 \mid \mathbf{x}]$	$. (1, . (2, \mathbf{x}))$
$[1, 2 \mid [3, 4]]$	same as $[1, 2, 3, 4]$

- Last in a list can be the symbol $|$ followed by a final term for the tail of the list
- Useful in patterns: $[1, 2 \mid \mathbf{x}]$ unifies with any list that starts with $1, 2$ and binds \mathbf{x} to the tail

$\begin{aligned} ?- [1, 2 \mid X] &= [1, 2, 3, 4, 5]. \\ \mathbf{x} &= [3, 4, 5]. \end{aligned}$
--

The **append** Predicate

```
?- append([1,2],[3,4],Z).  
Z = [1, 2, 3, 4].
```

- Predefined **append (X, Y, Z)** succeeds if and only if **Z** is the result of appending the list **Y** onto the end of the list **X**

Not Just A Function

```
?- append(X, [3, 4], [1, 2, 3, 4]) .  
x = [1, 2] .
```

- **append** can be used with any pattern of instantiation (that is, with variables in any positions)

Not Just A Function

```
?- append(X, Y, [1, 2, 3]) .  
X = [],  
Y = [1, 2, 3] ;  
X = [1],  
Y = [2, 3] ;  
X = [1, 2],  
Y = [3] ;  
X = [1, 2, 3],  
Y = [] ;  
false.
```

An Implementation

```
append([], B, B).  
append([Head|TailA], B, [Head|TailC]) :-  
    append(TailA, B, TailC).
```

Other Predefined List Predicates

Predicate	Description
member (X, Y)	Provable if the list Y contains the element X .
select (X, Y, Z)	Provable if the list Y contains the element X , and Z is the same as Y but with one instance of X removed.
nth0 (X, Y, Z)	Provable if X is an integer, Y is a list, and Z is the X th element of Y , counting from 0.
length (X, Y)	Provable if X is a list of length Y .

- All flexible, like **append**
- Queries can contain variables anywhere

Using **select**

```
?- select (2, [1, 2, 3], Z) .  
Z = [1, 3] ;  
false.
```

```
?- select (2, Y, [1, 3]) .  
Y = [2, 1, 3] ;  
Y = [1, 2, 3] ;  
Y = [1, 3, 2] ;  
false.
```

The **reverse** Predicate

```
?- reverse([1,2,3,4],Y).  
Y = [4, 3, 2, 1].
```

- Predefined **reverse (X, Y)** unifies **Y** with the reverse of the list **X**

An Implementation

```
reverse ([], []) .  
reverse ([Head|Tail], X) :-  
    reverse (Tail, Y) ,  
    append (Y, [Head], X) .
```

- Not an efficient way to reverse
- We'll see why, and a more efficient solution, in Chapter 21

Non-Terminating Queries

```
?- reverse(X, [1, 2, 3, 4]).  
X = [4, 3, 2, 1] ;
```

```
^CAction (h for help) ? abort  
% Execution Aborted  
?-
```

- Asking for another solution caused an infinite loop
- Hit Control-C to stop it, then *a* for abort
- **reverse** cannot be used as flexibly as **append**

Flexible and Inflexible

- Ideally, predicates should all be flexible like **append**
- They are more declarative, with fewer procedural quirks to consider
- But inflexible implementations are sometimes used, for efficiency or simplicity
- Another example is **sort...**

Example

```
?- sort([2,3,1,4],X).
```

```
X = [1, 2, 3, 4].
```

```
?- sort(X,[1,2,3,4]).
```

```
ERROR: Arguments are not sufficiently instantiated
```

- A fully flexible **sort** would also be able to unsort—find all permutations
- But it would not be as efficient for the more common task

The Anonymous Variable

- The variable `_` is an anonymous variable
- Every occurrence is bound independently of every other occurrence
- In effect, much like ML's `_`: it matches any term without introducing bindings

Example

```
tailof([_|A], A) .
```

- This **tailof**(**X**, **Y**) succeeds when **X** is a non-empty list and **Y** is the tail of that list
- Don't use this, even though it works:

```
tailof([Head|A], A) .
```

Dire Warning

```
append([], B, B).  
append([Head|TailA], B, [Head|TailC]) :-  
    append(TailA, B, TailC).
```

- Don't ignore warning message about singleton variables
- As in ML, it is bad style to introduce a variable you never use
- More importantly: *if you misspell a variable name, this is the only warning you will see*

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The **not** Predicate

```
?- member(1, [1, 2, 3]).  
true .  
  
?- not(member(4, [1, 2, 3])).  
false.
```

- For simple applications, it often works quite a bit like logical negation
- But it has an important procedural side...

Negation As Failure

- To prove **not (X)**, Prolog attempts to prove **X**
- **not (X)** succeeds if **X** fails
- The two faces again:
 - Declarative: **not (X) = $\neg X$**
 - Procedural: **not (X)** succeeds if **X** fails, fails if **X** succeeds, and runs forever if **X** runs forever

Example

```
sibling(X, Y) :-  
    not(X=Y),  
    parent(P, X),  
    parent(P, Y).
```

```
?- sibling(kim, kent).  
true .
```

```
?- sibling(kim, kim).  
false.
```

```
?- sibling(X, Y).  
false.
```

```
sibling(X, Y) :-  
    parent(P, X),  
    parent(P, Y),  
    not(X=Y).
```

```
?- sibling(X, Y).  
X = kim,  
Y = kent ;  
X = kent,  
Y = kim ;  
X = margaret,  
Y = jean ;  
X = jean,  
Y = margaret ;  
false.
```

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A Classic Riddle

- A man travels with wolf, goat and cabbage
- Wants to cross a river from west to east
- A rowboat is available, but only large enough for the man plus one possession
- Wolf eats goat if left alone together
- Goat eats cabbage if left alone together
- How can the man cross without loss?

Configurations

- Represent a configuration of this system as a list showing which bank each thing is on in this order: man, wolf, goat, cabbage
- Initial configuration: **[w, w, w, w]**
- If man crosses with wolf, new state is **[e, e, w, w]** – but then goat eats cabbage, so we can't go through that state
- Desired final state: **[e, e, e, e]**

Moves

- In each move, man crosses with at most one of his possessions
- We will represent these four moves with four atoms: **wolf**, **goat**, **cabbage**, **nothing**
- (Here, **nothing** indicates that the man crosses alone in the boat)

Moves Transform Configurations

- Each move transforms one configuration to another
- In Prolog, we will write this as a predicate:
move (Config, Move, NextConfig)
 - **Config** is a configuration (like **[w, w, w, w]**)
 - **Move** is a move (like **wolf**)
 - **NextConfig** is the resulting configuration (in this case, **[e, e, w, w]**)

The **move** Predicate

```
change (e, w) .
```

```
change (w, e) .
```

```
move ([X, X, Goat, Cabbage], wolf, [Y, Y, Goat, Cabbage]) :-  
    change (X, Y) .
```

```
move ([X, Wolf, X, Cabbage], goat, [Y, Wolf, Y, Cabbage]) :-  
    change (X, Y) .
```

```
move ([X, Wolf, Goat, X], cabbage, [Y, Wolf, Goat, Y]) :-  
    change (X, Y) .
```

```
move ([X, Wolf, Goat, C], nothing, [Y, Wolf, Goat, C]) :-  
    change (X, Y) .
```


Safe Configurations

- A configuration is safe if
 - At least one of the goat or the wolf is on the same side as the man, and
 - At least one of the goat or the cabbage is on the same side as the man

```
oneEq (X, X, _) .
```

```
oneEq (X, _, X) .
```

```
safe ([Man, Wolf, Goat, Cabbage]) :-
```

```
  oneEq (Man, Goat, Wolf) ,
```

```
  oneEq (Man, Goat, Cabbage) .
```

Solutions

- A solution is a starting configuration and a list of moves that takes you to **[e, e, e, e]**, where all the intermediate configurations are safe

```
solution([e, e, e, e], []).  
solution(Config, [Move|Rest]) :-  
  move(Config, Move, NextConfig),  
  safe(NextConfig),  
  solution(NextConfig, Rest).
```

Prolog Finds A Solution

```
?- length(X, 7), solution([w, w, w, w], X).  
X = [goat, nothing, wolf, goat, cabbage, nothing, goat] .
```

- Note: without the **length(X, 7)** restriction, Prolog would not find a solution
- It gets lost looking at possible solutions like **[goat, goat, goat, goat, goat...]**
- More about this in Chapter 20

What Prolog Is Good For

- The program specified a problem logically
- It did not say how to search for a solution to the problem – Prolog took it from there
- That's one kind of problem Prolog is especially good for
- More examples to come in Chapter 22