## Constructing a Parse Tree

## Initial Conditions

Before we begin we need to have two things:

1. A Grammar
2. A string we wish to parse

So then what is the procedure:

1. Use the grammar to produce a derivation resulting in the given string
2. Use the derivation to produce the parse tree

## Example 1:

We have the following grammar:

```
<S> ::= <round> <square> | <outer>
<round> ::= ( <round> ) | ( )
<square> ::= [ <square> ] | [ ]
<outer> ::= ( <outer> ] | ( <inner> ]
<inner> ::= ) <inner> [ | ) [
```

We will derive the following string: (())[[]]

## Example 1: Producing the derivation

A derivation is produce using the following steps:

1. Start with the start rule and select one of its options
2. We then continue to replace each rule until we reach a terminal working from left to right.
3. We repeat step 2 until all non-terminals are replaced by terminals, and we have produced the target string.

Goal: Derive ( ( ) ) [ [ ] ]

```
<S> => <round> <square>
    => ( <round> ) <square>
    => ( ( ) ) <square>
    => ( ( ) ) [ <square> ]
    => ( ( ) ) [ [ ] ] --> Finished
```


## Example 1: Producing the parse tree

```
<S> => <round> <square>
<round> <square>
=> ( <round> ) <square>
=> ( ( ) ) <square>
=> ( ( ) ) [ <square> ]
=> ( ( ) ) [ [ ] ]
```



Example 2

Grammar:

$$
\begin{aligned}
& \text { <S1> : : }=\text { <S1> + <S2> | <S2> } \\
& \text { <S2> ::= <S2> * <S3> | <S3> } \\
& \text { <S3> ::= ( <S1> ) | a | b | c }
\end{aligned}
$$

String: $\mathrm{a}+\mathrm{b} * \mathrm{c}$

Example 2: Derivation

Goal: Derive $\mathrm{a}+\mathrm{b} * \mathrm{c}$

$$
\begin{aligned}
& \text { <S1> => <S1> + <S2> } \\
& \text { => <S2> + <S2> } \\
& \text { => <S3> + <S2> } \\
& \Rightarrow \mathrm{a}+\langle\mathrm{S} 2> \\
& \text { => } \mathrm{a}+\langle\mathrm{S} 2\rangle *<\text { S3> } \\
& \Rightarrow \mathrm{a}+\langle\text { S3> * <S3> } \\
& \Rightarrow \mathrm{a}+\mathrm{b} *<\text { S3> } \\
& \Rightarrow \mathrm{a}+\mathrm{b} * \mathrm{c}
\end{aligned}
$$

Example 2: Parse tree

$$
\begin{aligned}
& \text { <S1> => <S1> + <S2> } \\
& \text { => <S2> + <S2> } \\
& \text { => <S3> + <S2> } \\
& =>a+\langle S 2> \\
& \Rightarrow \mathrm{a}+\langle\mathrm{S} 2\rangle *<\text { S3 }\rangle \\
& \text { => } \mathrm{a}+\langle\mathrm{S} 3>\text { * <S3> } \\
& =>\mathrm{a}+\mathrm{b} *<\mathrm{S} 3> \\
& \Rightarrow \mathrm{a}+\mathrm{b} * \mathrm{c}
\end{aligned}
$$

## We can also derive a string from the bottom up

Goal: Use Example 2 grammar to derive $\mathrm{a}+\mathrm{b} * \mathrm{c}$ In This case we start with the rightmost terminal and continue to replace with non-terminals until we reach the start rule.

$$
\begin{aligned}
& \Rightarrow \mathrm{a}+\mathrm{b} * \mathrm{c} \\
& \Rightarrow \mathrm{a}+\mathrm{b} *<\mathrm{S} 3> \\
& \text { => } \mathrm{a}+\langle\mathrm{S} 3>\text { * <S3> } \\
& \text { => } \mathrm{a}+\text { <S2> * <S3> } \\
& \Rightarrow \mathrm{a}+\langle\mathrm{S} 2\rangle \\
& \text { => <S3> + <S2> } \\
& \text { => <S2> + <S2> } \\
& \text { => <S1> + <S2> } \\
& \text { <S1> => <S1> + <S2> }
\end{aligned}
$$

We then build the parse tree starting from the bottom

Constructing Abstract Syntax Trees

## Parse Trees to ASTs

An Abstract Syntax Tree (AST) is a simplified form of a Parse Tree which is useful for interpreting/converting code from one language to another. Thus, it is useful for the compiling process.

An AST typically is a Binary Tree and requires that there are no non-terminal symbols left in the tree, and this then requires that we have the following:

- A Parse Tree
- A notion of traversing the tree (we will assume an In Order traversal)


## The Conversion Process



## Example

| <S1> |  |  |  | <S1> |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| / | \| | $\backslash$ |  | / | \| | $\backslash$ |  |
| <S1> | + | <S2> |  | a | + | <S2> |  |
| \| | / | I | 1 |  | / | I | $\backslash$ |
| <S2> | 1 | 1 | 1 |  | 1 | 1 | $\backslash$ |
| \| | 1 | 1 | $\backslash$ |  | 1 | I | $\backslash$ |
| <S3> | 1 | I | I |  | I | I | 1 |
| । | 1 | I | I |  | I | 1 | I |
| a | । | 1 | I |  | । | I | I |
|  | <S2> | * | <S3> |  | <S2> | * | <S3> |
|  | \| |  | I |  | 1 |  | । |
|  | <S3> |  | 1 |  | <S3> |  | I |
|  | I |  | 1 |  | I |  | I |
|  | b |  | 1 |  | b |  | 1 |
|  |  |  | c |  |  |  | c |

## Example

| <S1> |  |  |  | <S1> |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| / | । | $\backslash$ |  | / | \| | $\backslash$ |  |
| <S1> | + | <S2> |  | a | + | <S2> |  |
| \| | / | \| | $\backslash$ |  | / | \| | $\backslash$ |
| <S2> | 1 | 1 | 1 |  | b | I | $\backslash$ |
| \| | 1 | 1 | $\backslash$ |  |  | I | $\backslash$ |
| <S3> | 1 | 1 | 1 |  |  | I | 1 |
| \| | I | I | 1 |  |  | \\| | I |
| a | । | 1 | I |  |  | I | । |
|  | <S2> | * | <S3> |  |  | * | <S3> |
|  | \| |  | \| |  |  |  | 1 |
|  | <S3> |  | 1 |  |  |  | । |
|  | \| |  | 1 |  |  |  | I |
|  | b |  | 1 |  |  |  | 1 |
|  |  |  | c |  |  |  | c |

## Example

| <S1> |  |  |  | <S1> |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| / | । | $\backslash$ |  | / | \| | $\backslash$ |  |
| <S1> | + | <S2> |  | a | + | <S2> |  |
| \| | / | \| | $\backslash$ |  | / | \| | $\backslash$ |
| <S2> | 1 | 1 | 1 |  | b | * | $\backslash$ |
| \| | 1 | 1 | $\backslash$ |  |  |  | $\backslash$ |
| <S3> | 1 | 1 | 1 |  |  |  | 1 |
| \| | I | I | 1 |  |  |  | I |
| a | । | 1 | I |  |  |  | । |
|  | <S2> | * | <S3> |  |  |  | <S3> |
|  | \| |  | \| |  |  |  | \| |
|  | <S3> |  | 1 |  |  |  | । |
|  | \| |  | 1 |  |  |  | I |
|  | b |  | 1 |  |  |  | 1 |
|  |  |  | c |  |  |  | c |

## Example

| <S1> |  |  |  | <S1> |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | \| | $\backslash$ |  | / | I | $\backslash$ |  |
| <S1> | + | <S2> |  | a | + | <S2> |  |
| \| | / | I | 1 |  | 1 | 1 | $\backslash$ |
| <S2> | 1 | 1 | 1 |  | b | * | c |
| \| | 1 | 1 | $\backslash$ |  |  |  |  |
| <S3> | I | 1 | 1 |  |  |  |  |
| I | 1 | 1 | 1 |  |  |  |  |
| a | \| | 1 | 1 |  |  |  |  |
|  | <S2> | * | <S3> |  |  |  |  |
|  | \| |  | 1 |  |  |  |  |
|  | <S3> |  | 1 |  |  |  |  |
|  | \| |  | 1 |  |  |  |  |
|  | b |  | 1 |  |  |  |  |
|  |  |  | c |  |  |  |  |

## Example

| <S1> |  |  |  | <S1> |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| / | \| | $\backslash$ |  | 1 | I | $\backslash$ |
| <S1> | + | <S2> |  | a | $+$ | * |
| \| | 1 | I | 1 |  |  | / \} |
| <S2> | 1 | 1 | $\backslash$ |  |  | b c |
| \| | 1 | 1 | $\backslash$ |  |  |  |
| <S3> | 1 | I | 1 |  |  |  |
| \| | I | 1 | 1 |  |  |  |
| a | \| | 1 | 1 |  |  |  |
|  | <S2> | * | <S3> |  |  |  |
|  | \| |  | I |  |  |  |
|  | <S3> |  | 1 |  |  |  |
|  | I |  | 1 |  |  |  |
|  | b |  | 1 |  |  |  |
|  |  |  | c |  |  |  |

Example


## Example



## Example



## Example



## Example



## Example



## Example



## Example



## Example



## Example



